

# Minimum bisection of random 3-regular graphs

## Lyuben Lichev

Joint work with Dieter Mitsche

# Minimum bisection of random 3-regular graphs

- 1 Introduction
- 2 The lower bound
- 3 The upper bound

# The problem

- Graph  $G(V, E)$ ,  $|V|$  being even.
- Bisection of  $G$ : partition  $(V_1, V_2)$  of  $V$  with  $|V_1| = |V_2|$ .
- Size of a bisection  $(V_1, V_2)$ : number of edges crossing.
- *Bisection width of  $G$* ,  $\text{bw}(G)$ : minimum size of a bisection of  $G$ .
- Find  $\text{bw}(G)$ .

Wide range on applications:

- Applied mathematics.
- Computer science.
- Statistical physics.

A much studied problem.

- NP-complete for general graphs (Garey, Johnson, Stockmeyer, 1976).
- NP-complete for regular graphs (Bui, Chaudhuri, Leighton, Sisper, 1987).
- Lower bounds for random 3-regular graphs:
  - $n/11 \approx 0.0909n$  (Bollobás, 1984).
  - $10n/99 \approx 0.10101n$  (Kostochka and Melnikov, 1993).

Upper bounds are more abundant.

- $n/4 + o(n)$  for an arbitrary 3-regular graph (Kostochka and Melnikov, 1992).
- $n/6$  for an arbitrary large enough 3-regular graph (Monien and Preis, 2001).
- $0.1740n$  for a random 3-regular graph (Díaz, Do, Serna and Wormald, 2003). Weaker bound, but more general paper. Easier algorithmic approach.
- $0.16226n$  for a random 3-regular graph (Lyons, 2017).

# Our contribution

## Theorem (L. and Mitsche, 2020)

*The bisection width of the random 3-regular graph is between  $0.10329n$  and  $0.13983n$  asymptotically almost surely.*

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Main ideas of the proof:

Lower bound:

- precise structural characterisation of a minimum bisection,
- partially regroup these "special" bisections according to the graph they originate from,
- bound from above the number of graphs containing a "special" bisection of given size.



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Main ideas of the proof:

Upper bound:

- local limit theory ensures convergence to the 3-regular tree  $T_3$ ,
- deduce convergence of the " $\lambda$ -Gaussian wave" on  $G_3(n)$  to the " $\lambda$ -Gaussian wave" on  $T_3$ ,
- based on a well-chosen  $\lambda$ , find bisections of small size in  $G_3(n)$ .

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# The probability space

- The configuration model.
- A set of vertices  $V$ .
- A set of half-edges, attached to the vertices.
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Figure: G. Zamora-López, PhD thesis.

# The probability space

- The configuration model.
- A set of vertices  $V$ .
- A set of half-edges, attached to the vertices.
- Random perfect matching of the half-edges.
- This model produces a multigraph in general.
- Simple graph is sampled with probability, bounded away from 0 and 1 as  $n \rightarrow +\infty$ .

# Winning sets

- Fix 3-regular graph  $G$  and a cut  $(V_1, V_2)$ .
- Let  $S \subseteq V_i, i \in \{1, 2\}$ .
- For  $\ell \geq 0$ ,  $S$  is  $\ell$ -winning if

$$|(V_i \setminus S, V_{3-i} \cup S)| = |(V_1, V_2)| - \ell.$$

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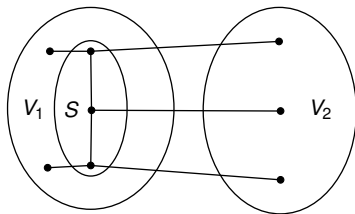


Figure: Here  $S$  is a 1-winning set.

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- Switching winning sets  $S, S'$  with  $|S| = |S'|$  and no edges between  $S$  and  $S'$  preserves the size of the parts and does not increase the cut.



# Improvement

- Fix 3-regular graph  $G$  and a bisection  $(V_1, V_2)$ .
- Let  $S_1 \subseteq V_1, S_2 \subseteq V_2$  such that

$$|(V_1, V_2)| > |(S_2 \cup (V_1 \setminus S_1), S_1 \cup (V_2 \setminus S_2))|.$$

- *Improvement of  $(V_1, V_2)$* : operation of switching  $S_1$  and  $S_2$ .

# Improvements and minimum bisections

- No minimum bisection admits an improvement!
- Switching winning sets  $S, S'$  with  $|S| = |S'|$ , not both indifferent (i.e. not changing the size of the cut) and without edges between  $S$  and  $S'$ , leads to an improvement.
- This gives information on the structure of the parts in a minimum bisection!

# Typical graphs

A 3-regular graph  $G$  is *typical* if:

- $\text{bw}(G) \geq n/10$ , and
- there are at most  $\log n$  vertices in cycles of length at most 20 in  $G$ .

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Almost all 3-regular graphs are typical.

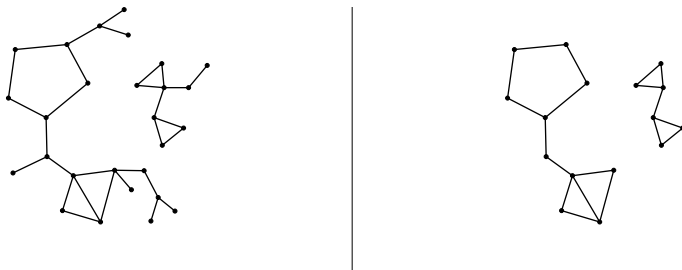
Goal: bound from above the number of minimum bisections of typical 3-regular graphs of size  $\beta n$ , for any  $\beta \geq 0.1$ .

# Step 1: 2-cores

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**Figure:** Left figure: a graph, right figure: its 2–core

# Step 1: 2-cores

## Lemma

*Let  $(V_1, V_2)$  be a minimum bisection of size  $\beta n$  of a typical 3-regular graph  $G$ . Then, for both  $i = 1, 2$ ,  $|G[V_i]| - |C_2(G[V_i])| \leq 4$ .*

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## Sketch of proof.

We argue by contradiction. We find many paths  $v_0 \dots v_k$  in  $G[V_{3-i}]$ , for some  $k \leq 4$ , such that

$$\deg_{G[V_{3-i}]}(v_0) \leq 2, \deg_{G[V_{3-i}]}(v_k) \leq 2.$$

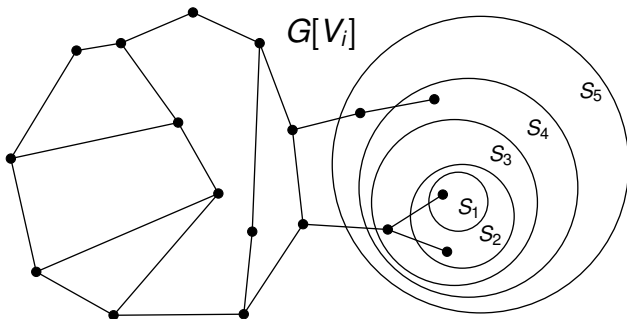
Exchanging one of them with the set  $G[V_i] \setminus C_2(G[V_i])$  would lead to an improvement. □



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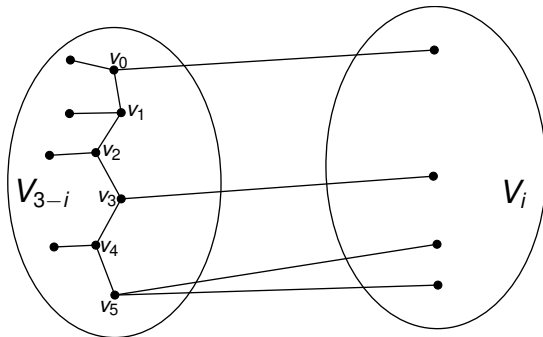
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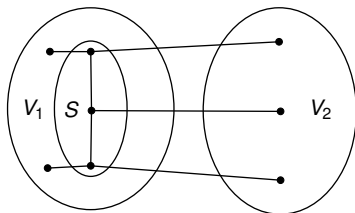
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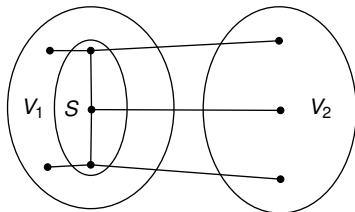
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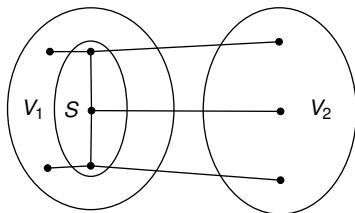
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Two cases:

- There are at most  $o(n)$  pairs of neighbours of degree two in  $G[V_2]$ .
- There is at most one path like  $S$  in  $G[V_2]$  (parity constraints).

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## Observation

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First moment method: only a negligible proportion of all 3-regular graphs admit such a bisection of size  $\beta n \leq 0.1069n$ .



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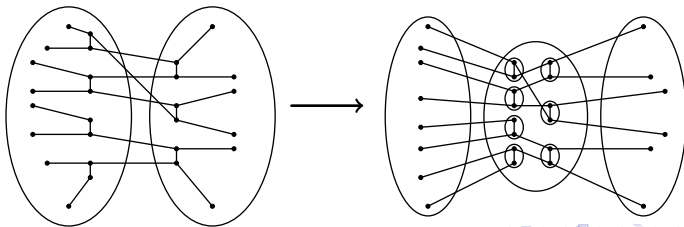
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New idea: regroup "special" bisections of second type according to the graph these come from.



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# Local limit theory

A *uniformly rooted graph*  $(G, \rho)$  is a finite graph with a root, chosen uniformly at random. A sequence of uniformly rooted graphs  $(G_n, \rho_n)_{n \geq 1}$  *converges locally* if for every  $r \geq 1$ ,

$$\lim_{n \rightarrow +\infty} \mathbb{P}(B_{G_n}(\rho_n, r) = (G, \rho))$$

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The limit measure of  $(B_{G_n}(\rho_n, r))_{n \geq 1}$  is often abusively identified with a graph  $G_\infty$ . Example: the neighbourhood of a typical vertex in  $G_3(n)$  contains no cycle of constant length. Therefore,  $(G_3(n))_{n \geq 1}$  converges to the 3-regular tree  $T_3$ .

# Factor of iid process

Given a graph  $G$ , assign a family  $(Z_v)_{v \in V(G)}$  of iid standard normal random variables to the vertices of  $G$ .

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A *factor of iid process* on  $G$  is a family of random variables  $(X_v)_{v \in G}$  such that

- for every  $v \in V(G)$ ,  $X_v$  is a measurable function of  $(Z_v)_{v \in V(G)}$ , and
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$(X_v)_{v \in V(G)}$  is a *linear factor of iid process* on  $G$  if there exist  $\alpha_0, \alpha_1, \dots$  such that  $X_v = \sum_{u \in V(G)} \alpha_{d_G(v,u)} Z_u$  for all  $v \in V(G)$ .

# Gaussian processes

A collection of random variables  $(Y_v)_{v \in V(G)}$  a *Gaussian process* on  $G$  if  $(Y_v)$  are jointly Gaussian and  $Y_v$  is centered for every  $v \in V(G)$ .

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**Theorem (Czóka, Gerencsér, Harangi, Virág, 2008)**

*For any real number  $\lambda$  with  $|\lambda| \leq 3$  there exists a non-trivial invariant Gaussian process  $(Y_v)_{v \in V(T_3)}$  on  $T_3$  such that (with probability 1) for every vertex  $v$*

$$\sum_{u \in N(v)} Y_u = \lambda Y_v.$$

*Moreover, the joint distribution of such a process is unique under the additional condition that, for every  $v \in V(T_3)$ , the variance of  $Y_v$  is 1. We refer to this process as the Gaussian wave function corresponding to the eigenvalue  $\lambda$ .*

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Let  $(G_n)_{n \geq 1}$  converge locally to  $T_3$ .

## Theorem (Czóka, Gerencsér, Harangi, Virág, 2008)

*For any real number  $\lambda$  with  $|\lambda| \leq 2\sqrt{2}$  there exists a sequence of linear factors of iid processes  $(X_v^{(n)})_{v \in V(G_n)}$  that converges in distribution to the Gaussian wave function  $(Y_v)_{v \in V(T_3)}$  corresponding to the eigenvalue  $\lambda$ .*

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Idea (Lyons, 2017): Form a cut of the vertices of  $G_3(n)$  according to the signs of  $(X_v)_{v \in V(G_3(n))}$ .

# Building upon Lyons' idea

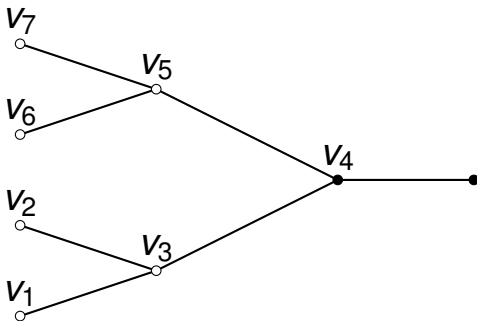
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# Exact results

- A study of the above structure leads to an improved upper bound of  $0.1398n$  on the bisection width.
- Studying longer chains of repeating black and white vertices is out of reach for us due to numerical difficulties. Monte-Carlo simulations allow us to make a non-rigorous prediction of an upper bound of  $0.1303n$ .



Thank you for your attention!

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